

LA-UR-17-24963

Approved for public release; distribution is unlimited.

Title: Numerically Solving the Heat Equation

Author(s): Garrett, Charles Kristopher

Intended for: 2017 LANL Parallel Computing Summer Research Internship lecture

Issued: 2017-06-20



Los Alamos National Laboratory LA-UR-17-24453

Numerically Solving the Heat Equation



Kris Garrett

June 2017













The Heat Equation

1D Equation

$$\frac{\partial u}{\partial t} = \frac{\partial^2 u}{\partial x^2}$$

2D Equation

$$\frac{\partial u}{\partial t} = \frac{\partial^2 u}{\partial x^2} + \frac{\partial^2 u}{\partial y^2}$$

Boundary and Initial Conditions in 1D

$$u(a,t) = g_a(t)$$

$$u(b,t) = g_b(t)$$

$$u(x,0) = u_0(x)$$

Will concentrate on the 1D equation for this presentation



Discretization in Space/Time

Exact solution

$$u_i^n \approx u(x_i, t^n)$$

 $u_i^n pprox u(x_i, t^n)$ Approximate solution in x and t

Common discretization of second derivative

$$\frac{\partial^2 u}{\partial x^2} = \frac{u(x - \Delta x) - 2u(x) + u(x + \Delta x)}{\Delta x^2} + O(\Delta x^2)$$

Means a quantity less than C Δx^2 where C is problem dependent but not discretization dependent



How to Get the Approximation

$$\frac{\partial^2 u}{\partial x^2} = \frac{u(x - \Delta x) - 2u(x) + u(x + \Delta x)}{\Delta x^2} + O(\Delta x^2)$$

Taylor Series

$$u(x + \Delta x) = u(x) + \Delta x u'(x) + \frac{\Delta x^2}{2} u''(x) + \frac{\Delta x^3}{6} u'''(x) + O(\Delta x^4)$$

$$u(x - \Delta x) = u(x) - \Delta x u'(x) + \frac{\Delta x^2}{2} u''(x) - \frac{\Delta x^3}{6} u'''(x) + O(\Delta x^4)$$

Add these, subtract 2u(x), and divide by Δx^2



Discretize in Time

Explicit Euler (1st Order)

$$u_i^{n+1} = u_i^n + \frac{\Delta t}{\Delta x^2} \left(u_{i-1}^n - 2u_i^n + u_{i+1}^n \right)$$

Implicit Euler (1st Order)

$$u_i^{n+1} = u_i^n + \frac{\Delta t}{\Delta x^2} \left(u_{i-1}^{n+1} - 2u_i^{n+1} + u_{i+1}^{n+1} \right)$$

Crank-Nicholson (2nd Order)

$$u_i^{n+1} = u_i^n + \frac{\Delta t}{2\Delta x^2} \left(u_{i-1}^{n+1} - 2u_i^{n+1} + u_{i+1}^{n+1} + u_{i-1}^n - 2u_i^n + u_{i+1}^n \right)$$



Discretize in Time

- "Method of Lines" approach
 - Discretized in all variables except time
 - Then discretize in time
- Local truncation error for time integration: O(Δt^{p+1})
 - Error from one time step
- Global error for time integration: O(Δt^p)
 - Error after all time steps
 - Order reduces by 1 by accumulating all time step errors



Two Things Always Needed

- Stability
- Consistency

Lax Equivalence Theorem for Linear PDEs

Convergence = Stable + Consistent



What is Consistent?

- Discretization converges to original PDE
- Example with implicit/explicit Euler steps

$$u_i^1 = u(x_i, t^1) + O(\Delta x^2) + O(\Delta t^2)$$

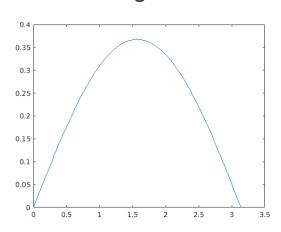


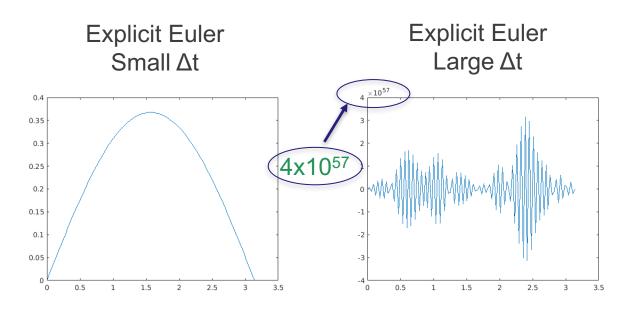
Importance of Stability

Initial Condition: $u_0(x) = \sin(x)$

Boundary Conditions: $u(0) = u(\pi) = 0$

Implicit Euler Large ∆t







Importance of Stability

- For stability
 - Explicit time integration: always subject to time step restriction
 - Typically the time step restriction is based on the size of Δx
 - Implicit time integration:
 - Sometimes no time step restriction
 - Usually less restriction on time step
 - One technique to calculate time step restriction is Von Neumann Analysis
- Remember: accuracy requires a time step restriction as well



Steps

- Apply Discrete Fourier Transform at time step n
- Apply time step to one Fourier mode
- See what conditions cause increase in size of Fourier mode
- Stability requires all initial conditions are damped
 - Keeps discretization and round off errors from increasing exponentially



Step 1: Apply Discrete Fourier Transform at time step n

$$u_j^n = \sum_{k=0}^{N-1} \hat{u}_k^n e^{-i(2\pi j/N)k}$$

$$\theta_j$$



Step 2: Apply time step to one Fourier mode

$$\hat{u}_k^{n+1} e^{-i\theta_j k} = \hat{u}_k^n e^{-i\theta_j k} + \frac{\Delta t}{\Delta x^2} \left(\hat{u}_k^n e^{-i\theta_{j-1} k} - 2\hat{u}_k^n e^{-i\theta_j k} + \hat{u}_k^n e^{-i\theta_{j+1} k} \right)$$

... do some algebra ...

$$\hat{u}_k^{n+1} = \left[1 + \frac{\Delta t}{\Delta x^2} \left(e^{i(2\pi/N)k} - 2 + e^{-i(2\pi/N)k}\right)\right] \hat{u}_k^n$$

Amplification Factor



Step 3: See what conditions cause increase in size of Fourier mode

Want value in brackets < 1
Maximum absolute value in brackets occurs for k = N/2

$$\left| 1 - 4 \frac{\Delta t}{\Delta x^2} \right| < 1$$

Leads to the stability requirement

$$\Delta t < \frac{1}{2} \Delta x^2$$



Implicit Methods

- Von Neumann analysis for implicit Euler and Crank-Nicholson
 - No time step restriction for stability
- But you have to solve a linear system
 - Takes more time to solve than explicit method



Implicit Methods

Linear System for implicit Euler or Crank-Nicholson

$$\begin{pmatrix} 1 + 2r & -r & & & \\ -r & 1 + 2r & -r & & & \\ & \ddots & \ddots & \ddots & & \\ & & -r & 1 + 2r & -r \\ & & & -r & 1 + 2r \end{pmatrix} u^{n+1} = b^n$$



Implicit Methods

Nice properties of linear system

- Strictly diagonally dominant
 - Gershgorin circle theorem implies system is positive definite
- Symmetric system
 - Implies system is diagonalizable (basis of eigenvectors)
 - All eigenvalues are real

If Δt is approximately Δx, then r is very big

- As Δx goes to zero, system approaches weakly diagonally dominant
- Harder for many iterative methods to converge



Solving Linear Systems

- Dense linear algebra: O(N³) flops
 - Gaussian Elimination (Lapack, Scalapack)
 - Can exploit banded nature of linear system
 - Can exploit sparse nature of linear system (SuperLU)
- Sparse linear algebra: $O(I N^2)$ flops I = Iterations
 - Classical: Jacobi, Gauss-Seidel, SOR (Usually hand coded)
 - Krylov: CG, GMRES, etc (Petsc, Trilinos)
 - Multigrid, Algebraic Multigrid (Hypre, Petsc, Trilinos)
 - Never stores the zeros of the matrix



- A = LU
 - L lower triangular, U upper triangular
 - Factorizing takes the most time: O(N³) flops
 - Solve Ax = b via LUx = b
 - Each triangular solve takes: O(N²) flops
 - Can reuse L and U for later solves
 - Really use A = PLU (Gaussian elimination with pivoting)
- For symmetric, positive definite: use Cholesky factorization
 - $A = L L^T$
 - No pivoting needed



- Uses BLAS (Basic Linear Algebra Subroutines)
 - Highly optimized: MKL, ACML, cuBLAS, ATLAS, OpenBLAS
 - Implements for example:
 - Matrix matrix multiply
 - Matrix vector multiply
 - Dot product of vectors



- LAPACK and BLAS originally FORTRAN libraries
 - CBLAS and LAPACKE for C interface
 - Can link to Fortran library from C/C++
- LAPACK library variants
 - ScaLAPACK MPI version
 - MAGMA GPU version
 - SuperLU unsymmetric, sparse systems

Parallel Computing SRI No constitute No constitute

- Other classical decompositions
 - QR decomposition: A = QR
 - Q is an orthogonal matrix
 - R is an upper triangular matrix
 - Used for least squares problems
 - Eigendecomposition
 - A = QDQ^T for symmetric problems
 - A = QTQ* for nonsymmetric problems
 - Q is orthogonal or hermitian
 - D is diagonal, real
 - T is triangular



Classical Sparse Linear Solvers

- Write A = N-M. Solve (N-M)x = b
 - Iterate $Nx^{k+1} = Mx^k + b$
 - Converges if and only if $\rho(N^{-1}M) < 1$ (all eigenvalue magnitudes < 1)
 - Jacobi: N is the diagonal of A
 - Gauss Seidel: N is the upper or lower triangular part of A
- Easy to code but converges slowly



Classical Sparse Linear Solvers

Linear System
$$(1+2r)x_i - rx_{i-1} - rx_{i+1} = b_i$$

Jacobi Method
$$(1+2r)x_i^{k+1} = rx_{i-1}^k + rx_{i+1}^k + b_i$$

Gauss Seidel Method

$$(1+2r)x_i^{k+1} - rx_{i-1}^{k+1} = rx_{i+1}^k + b_i$$



Krylov Linear Solvers

- Gets the "best answer" from a Krylov subspace K^k(A,b) = {b, Ab, A²b, ..., A^{k-1}b}
 - CG (Conjugate Gradient) used for symmetric, positive definite systems
 - Three vector recurrence relation
 - GMRES (Generalized Minimal Residual) used for nonsymmetric systems
 - Must hold all vectors in Krylov space
 - Actually use GMRES(m): restart after m steps to reduce memory required
 - Guaranteed convergence for positive definite systems
 - Note: other Krylov spaces are used for some Krylov solvers



Krylov Linear Solvers

```
% Matlab version of CG from Wikipedia
function [x] = conjgrad(A, b, x)
    r = b - A * x;
    rsold = r' * r;
    for i = 1:length(b)
        Ap = A * p;
        alpha = rsold / (p' * Ap);
        x = x + alpha * p;
        r = r - alpha * Ap;
        rsnew = r' * r; ---
        if sqrt(rsnew) < 1e-10
            break;
        end
        p = r + (rsnew / rsold) * p;
        rsold = rsnew;
    end
end
```

Notice only 3 extra vectors of memory required

Also need to perform dot products. Can hurt parallel performance



Multigrid Linear Solver

Uses the classical solvers

- These solvers converge quickly for certain discrete Fourier modes
- When grid size changes, other Fourier modes converge quickly
- Solves problem on grid sizes: h, 2h, 4h, 8h, etc.
- Generally fastest solver for diffusion type equations



Iterative Solvers

- Many iterative solvers need a preconditioner
 - Ax = b
 - PAx = Pb (Left preconditioner)
 - PA should require less iterations
 - P should be easily invertible
 - There are also right and symmetric preconditioners
 - Very problem dependent

Parallel Computing SRI



Other Important Linear Solver Topics

Norm of vector ||v||

$$||v||_1 = |v_1| + |v_2| + \dots + |v_n|$$

$$||v||_2 = (v_1^2 + v_2^2 + \dots + v_n^2)^{1/2}$$

$$||v||_{\infty} = \max\{|v_1|, |v_2|, \dots, |v_n|\}$$

Induced norm of matrix (max matrix stretches a vector)

$$||A|| = \max_{||v||=1} ||Av||$$

Parallel LSTI



Other Important Linear Solver Topics

Condition number of a matrix

$$\kappa(A) = ||A^{-1}|| \cdot ||A||$$

 If you solve Ax = b+e, the relative error in solution compared to the relative error in RHS is

$$\frac{||A^{-1}e||}{||A^{-1}b||} \le \kappa(A) \frac{||e||}{||b||}$$

- This is for exact arithmetic
- This shows the error in solution given error in data



Ensuring Correctness

- You must run tests to ensure a correct answer and correct implementation
- Verification Tests: Make sure you are actually solving the heat equation
 - Use known analytical solutions: sin(x) e^{-t}
 - Method of manufactured solution: used when there is a known source.
 - Make up a solution and determine the source.
 - Put that source into your solver.



Ensuring Correctness

- Convergence Tests: Make sure numerical implementation is correct
 - It is common to code incorrectly and get first order convergence of a higher order method
- Unit Tests: Test code in every file
 - Very useful for large projects
 - When you find a bug, add a test to reproduce it
 - Makes pinpointing errors easier
 - Some projects require that every branch in code is tested

The End